

MODULARITY AND DISTRIBUTIVITY OF THE LATTICE  
OF  $\Sigma$ -CLOSED SUBSETS OF AN ALGEBRAIC STRUCTURE

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*Summary.* Let  $\mathcal{A} = (A, F, R)$  be an algebraic structure of type  $\tau$  and  $\Sigma$  a set of open formulas of the first order language  $L(\tau)$ . The set  $C_\Sigma(\mathcal{A})$  of all subsets of  $A$  closed under  $\Sigma$  forms the so called lattice of  $\Sigma$ -closed subsets of  $\mathcal{A}$ . We prove various sufficient conditions under which the lattice  $C_\Sigma(\mathcal{A})$  is modular or distributive.

*Keywords:* algebraic structure, closure system,  $\Sigma$ -closed subset, modular lattice, distributive lattice, convex subset

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Modularity and distributivity of subalgebra lattices was investigated by T. Evans and B. Ganter in [4] and by the first author in [1]. However, we can study much more general lattices of closed subsets of an algebra or a relational structure. For convex sublattices of a given lattice this was done by V. I. Marmazajev [6], for convex subsets of monounary algebras or ordered sets see [5] or [3], respectively. A general approach for these considerations was developed by the authors in [2]. By using it, we can state sufficient (and in some cases also necessary) conditions under which a lattice of all  $\Sigma$ -closed subsets of a given algebraic structure is modular or even distributive.

First we recall some concepts. By a *type* we mean a pair of sequences  $\tau = \langle \{n_i; i \in I\}, \{m_j; j \in J\} \rangle$  where  $n_i, m_j$  are non-negative integers. An *algebraic structure* or briefly a *structure* of type  $\tau$  is a triplet  $\mathcal{A} = (A, F, R)$ , where  $A \neq \emptyset$  is a set and  $F = \{f_i; i \in I\}$ ,  $R = \{\varrho_j; j \in J\}$  such that for each  $i \in I$ ,  $f_i$  is an  $n_i$ -ary *operation* on  $A$  and for each  $j \in J$ ,  $\varrho_j$  is an  $m_j$ -ary *relation* on  $A$ . Denote by  $L(\tau)$  the first order language containing operational and relational symbols of type  $\tau$ . If  $R = \emptyset$ , the structure  $(A, F, \emptyset)$  is denoted briefly by  $(A, F)$  and is called *an algebra*. If  $F = \emptyset$  then  $(A, \emptyset, R)$  is denoted by  $(A, R)$  and called *a relational system*; this system

$(A, R)$  is called *binary* if each  $\varrho_j \in R$  is binary. A binary relational system  $(A, R)$  is said to be *antisymmetrical* if each  $\varrho_j \in R$  is an antisymmetrical relation. A binary relational system  $(A, R)$  is called an *ordered* (or *quasiordered*) *set* if  $R = \{\varrho_1\}$  where  $\varrho_1$  is an *order* on  $A$  (or a reflexive and transitive relation, the so called *quasiorder*, respectively).

Let  $\Gamma$  be an index set and for each  $\gamma \in \Gamma$  let  $G_\gamma(x_1, \dots, x_{k_\gamma}, y_1, \dots, y_{s_\gamma}, z, f_i)$  be an open formula of a language  $L(\tau)$  containing individual variables  $x_1, \dots, x_{k_\gamma}, y_1, \dots, y_{s_\gamma}, z$  and a symbol  $f_i$  of  $n_i$ -ary term operation. Analogously, let  $\Lambda$  be an index set and for each  $\lambda \in \Lambda$  let  $G_\lambda(x_1, \dots, x_{k_\lambda}, y_1, \dots, y_{s_\lambda}, z, \varrho_j)$  be an open formula of the language  $L(\tau)$  containing individual variables  $x_1, \dots, x_{k_\lambda}, y_1, \dots, y_{s_\lambda}, z$  and a symbol  $\varrho_j$  of  $m_j$ -ary relation of type  $\tau$ . Put  $\Sigma = \{G_\gamma; \gamma \in \Gamma\} \cup \{G_\lambda; \lambda \in \Lambda\}$ .

**Definition 1.** A subset  $B$  of an algebraic structure  $\mathcal{A} = (A, F, R)$  is called  $\Sigma$ -closed if for every  $\gamma \in \Gamma, \lambda \in \Lambda$  and  $a_1, \dots, a_{k_\gamma}, a'_1, \dots, a'_{k_\lambda} \in B$  and  $b_1, \dots, b_{s_\gamma}, b'_1, \dots, b'_{s_\lambda}, c, c' \in A$ , we have  $c \in B$  or  $c' \in B$  provided  $G_\gamma(a_1, \dots, a_{k_\gamma}, b_1, \dots, b_{s_\gamma}, c, f_i)$  or  $G_\lambda(a'_1, \dots, a'_{k_\lambda}, b'_1, \dots, b'_{s_\lambda}, c', \varrho_j)$  are satisfied in  $\mathcal{A}$ . Denote by  $C_\Sigma(\mathcal{A})$  the set of all  $\Sigma$ -closed subsets of  $\mathcal{A}$ .

As was proved in [2], the set  $C_\Sigma(\mathcal{A})$  of all  $\Sigma$ -closed subsets of a structure  $\mathcal{A} = (A, F, R)$  is a complete lattice with respect to set inclusion with the greatest element  $A$ . In what follows we will study modularity and distributivity of  $C_\Sigma(\mathcal{A})$  depending on the properties of  $\mathcal{A}$ . For any given structure  $\mathcal{A}$  we will suppose that the set of formulas  $\Sigma$  is determined. For a given subset  $M \subseteq A$  we denote by  $C_{\mathcal{A}}(M)$  the least  $\Sigma$ -closed subset of  $\mathcal{A}$  containing  $M$ ; we say that  $C_{\mathcal{A}}(M)$  is *generated* by  $M$ . If  $M$  is a finite subset, say  $M = \{a_1, \dots, a_k\}$ , we will write  $C_{\mathcal{A}}(a_1, \dots, a_k)$  for  $C_{\mathcal{A}}(M)$ .

If the set  $\Sigma$  is implicitly known, we will use on the lattice  $C_\Sigma(\mathcal{A})$  to specify the closure system. In some more familiar examples of  $C_\Sigma(\mathcal{A})$  we will use the common name and notation:

(1) If  $\mathcal{A} = (A, F)$  is an algebra,  $F = \{f_i: i \in I\}$  and  $\Sigma = \{G_i: i \in I\}$  where  $G_i(x_1, \dots, x_{n_i}, z, f_i)$  is the formula  $(f_i(x_1, \dots, x_{n_i}) = z)$ , then  $\Sigma$ -closed subsets of  $\mathcal{A}$  are subalgebras of  $\mathcal{A}$  and  $\emptyset$ , and  $C_\Sigma(\mathcal{A}) = \text{Sub } \mathcal{A}$ .

(2) If  $\mathcal{L} = (L, \{\vee, \wedge\})$  is a lattice,  $\Sigma = \{G_1, G_2\}$  where  $G_1$  is the formula  $(x_1 \vee x_2 = z)$  and  $G_2$  is the formula  $(x_1 \wedge y_1, z)$ , then the  $\Sigma$ -closed subsets of  $\mathcal{L}$  are lattice ideals, i.e.  $C_\Sigma(\mathcal{L}) = \text{Id } \mathcal{L}$ .

(3) If  $\mathcal{R} = (A, R)$  is a binary relational system with  $R = \{\varrho_j; j \in J\}$  and  $\Sigma = \{G_j: j \in J\}$  where for each  $j \in J$  we have

$$G_j \text{ is the formula } (x_1 \varrho_j z \text{ and } z \varrho_j x_2),$$

then the  $\Sigma$ -closed subsets of  $\mathcal{R}$  are the so called *convex subsets* and  $C_\Sigma(\mathcal{R})$  will be denoted by  $\text{Conv } \mathcal{R}$ .

In particular, if  $\mathcal{S} = (S, \leq)$  is an ordered set then  $\Sigma = \{G\}$  where  $G$  is the formula  $(x_1 \leq z \leq x_2)$ . Thus  $\Sigma$ -closed subsets of  $\mathcal{S}$  are exactly the convex subsets of  $\mathcal{S}$  in the usual sense.

(4) If  $\mathcal{G} = (G, \cdot, ^{-1}, e)$  is a group and  $\Sigma = \{G_1, G_2, G_3, G_4\}$ , where  $G_1(x_1, x_2, z, \cdot)$  is the formula  $(x_1 \cdot x_2 = z)$ ,  $G_2(x_1, z, ^{-1})$  is the formula  $(x_1^{-1} = z)$ ,  $G_3(z, e)$  is the formula  $(e = z)$  and  $G_4(x_1, y_1, z, p)$  is the formula  $(p(x_1, y_1) = z)$  where  $p(x_1, y_1)$  is the term operation  $y_1 x_1 y_1^{-1}$ , then  $C_\Sigma(\mathcal{G})$  is the lattice of all normal subgroups of  $\mathcal{G}$ . It will be denoted simply by  $N(\mathcal{G})$ .

In what follows we denote join in  $C_\Sigma(\mathcal{A})$  by  $\vee$ , meet evidently coincides with set intersection.

**Theorem 1.** *Let  $\mathcal{A} = (A, F, R)$  be an algebraic structure with the system  $C_\Sigma(\mathcal{A})$  of  $\Sigma$ -closed subsets satisfying*

(i) *for each  $X, Y \in C_\Sigma(\mathcal{A})$ ,  $\emptyset \neq X \neq Y \neq \emptyset$  we have  $a \in X \vee Y$  if and only if there exist  $x \in X, y \in Y$  with  $a \in C_{\mathcal{A}}(x, y)$ ;*

(ii) *for each  $x, y \in A$ , if  $a \in C_{\mathcal{A}}(x, y)$  and  $C_{\mathcal{A}}(a) \neq C_{\mathcal{A}}(x)$  then  $y \in C_{\mathcal{A}}(x, a)$ . Then the lattice  $(C_\Sigma(\mathcal{A}), \subseteq)$  is modular.*

*Proof.* Suppose  $X, Y, Z \in C_\Sigma(\mathcal{A})$  and  $X \subseteq Z$ . If either  $X = \emptyset$  or  $Y = \emptyset$  the proof is trivial. Also for  $X = Y$  we easily obtain the modularity law. Hence, consider  $\emptyset \neq X \neq Y \neq \emptyset$ . Suppose  $a \in (X \vee Y) \cap Z$ . Then  $a \in Z$  and  $a \in X \vee Y$ . By (i), there exist  $x \in X, y \in Y$  such that  $a \in C_{\mathcal{A}}(x, y)$ .

If  $C_{\mathcal{A}}(a) = C_{\mathcal{A}}(x)$  then  $a \in C_{\mathcal{A}}(a) = C_{\mathcal{A}}(x) \subseteq X \vee (Y \cap Z)$ .

If  $C_{\mathcal{A}}(a) \neq C_{\mathcal{A}}(x)$ , then we have  $y \in C_{\mathcal{A}}(x, a)$  by (ii).

However,  $x \in X \subseteq Z, a \in Z$  thus also  $y \in C_{\mathcal{A}}(x, a) \subseteq Z$ . Hence  $y \in Y \cap Z$  and  $a \in C_{\mathcal{A}}(x, y) \subseteq X \vee (Y \cap Z)$ , which proves modularity of  $C_\Sigma(\mathcal{A})$ .  $\square$

**Lemma 1.** *Let  $\mathcal{A} = (A, \{\rho\})$  be a binary relational system with only one transitive binary relation and  $C_\Sigma(\mathcal{A}) = \text{Conv } \mathcal{A}$ . Then  $C_\Sigma(\mathcal{A})$  satisfies (i) of Theorem 1.*

*Proof.* The condition (i) of Theorem 1 is equivalent to the following one:

$$C_{\mathcal{A}}(X) = \bigcup \{C_{\mathcal{A}}(x_1, x_2); x_1, x_2 \in X\} \quad \text{for each } X \subseteq A.$$

For  $X, Y \subseteq A$  put  $C^0(X, Y) = X \cup Y, C(X, Y) = C^1(X, Y) = \{a \in A; u \rho a \rho v \text{ for some } u, v \in X \cup Y\}$  and  $C^{n+1}(X, Y) = C(C^n(X, Y))$ , where  $n \in N_0$  (non-negative integer). Evidently,  $C_{\mathcal{A}}(X, Y) = \bigcup \{C^n(X, Y); n \in N_0\}$ . Now, we can prove the following statement by induction on  $n$ : "If  $a \in C^n(X, Y)$ , then there exist  $u, v \in X \cup Y$  such that  $u \rho a \rho v$ ."

1) For  $n = 1$  it is a trivial.

2) Suppose that it is valid for all  $k \leq n$  and we prove it for  $n + 1$ . Let  $a \in C^{n+1}(X, Y)$ , i.e.  $\alpha \varrho a \varrho \beta$  for some  $\alpha, \beta \in C^n(X, Y)$ . Clearly, we have the following possibilities:

a)  $\alpha \in [x_1, y_1], \beta \in [x_1, y_2]$ ;

b)  $\alpha \in [x_1, y_1], \beta \in [y_2, x_2]$ ;

c)  $\alpha \in [y_1, x_1], \beta \in [x_2, y_2]$ ;

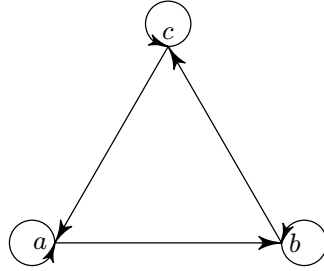
d)  $\alpha \in [y_1, x_1], \beta \in [y_2, x_2]$ , etc., where  $x_1, x_2 \in X$  and  $y_1, y_2 \in Y$ .

ad a) If  $\alpha \in [x_1, y_1], \beta \in [x_1, y_2]$ , then  $x_1 \varrho \alpha a \varrho \beta \varrho y_2$  and  $a \in [x_1, y_2]$  by transitivity, i.e. the statement is valid.

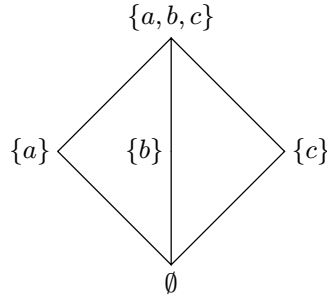
ad b)  $\alpha \in [x_1, y_1], \beta \in [y_2, x_2]$  imply  $x_1 \varrho \alpha a \varrho \beta \varrho x_2$ , i.e.  $a \in [x_1, x_2]$  and  $a \in X$ .

Similarly we can easily check the other possibilities.  $\square$

**Example 1.** Let  $\mathcal{A} = (\{a, b, c\}, \{\varrho\})$  be a binary relational system with the following diagram of  $\varrho$ :



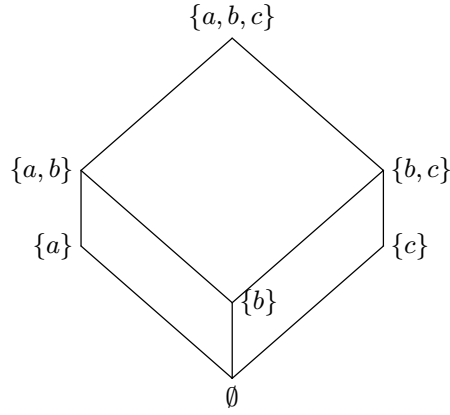
and  $C_\Sigma(\mathcal{A}) = \text{Conv } \mathcal{A}$ . We can easily check (i) and (ii) of Theorem 1, thus  $C_\Sigma(\mathcal{A})$  is modular. We can visualize the diagram of  $C_\Sigma(\mathcal{A})$  in Fig. 2 below:



We can see that it is isomorphic to  $M_3$ , hence  $C_\Sigma(\mathcal{A})$  is not distributive.

**Example 2.** Let  $\mathcal{A} = (\{a, b, c\}, \leq)$  be an ordered set which is a chain:  $a < b < c$ , and let  $C_\Sigma(\mathcal{A}) = \text{Conv } \mathcal{A}$ . Then it does not satisfy (ii) of Theorem 1 since  $b \in C_{\mathcal{A}}(a, c)$ ,  $C_{\mathcal{A}}(b) \neq C_{\mathcal{A}}(a)$  but  $c \notin C_{\mathcal{A}}(a, b) = \{a, b\}$ . The diagram of  $C_\Sigma(\mathcal{A})$  is

shown in Fig. 3:



We can see that  $C_{\Sigma}(\mathcal{A})$  is not modular. We are going to show that for some algebraic structures the condition (ii) is really equivalent to modularity of  $C_{\Sigma}(\mathcal{A})$ .

Recall from [2] that an algebraic system  $\mathcal{A} = (A, F, R)$  is  $\Sigma$ -separable if we have  $C_{\mathcal{A}}(x) = \{x\}$  for any  $x \in A$ .

**Theorem 2.** *Let  $\mathcal{A} = (A, F, R)$  be a  $\Sigma$ -separable algebraic structure satisfying (i) of Theorem 1. The following conditions are equivalent:*

- (a) *the lattice  $C_{\Sigma}(\mathcal{A})$  is modular;*
- (b) *for each  $x, y \in A$ , if  $a \in C_{\mathcal{A}}(x, y)$  for  $a \neq x$  then  $y \in C_{\mathcal{A}}(x, a)$ .*

*Proof.* Since  $\mathcal{A}$  is  $\Sigma$ -separable and  $\mathcal{A}$  satisfies (i), we obtain (b)  $\Rightarrow$  (a) directly by Theorem 1. Prove (a)  $\Rightarrow$  (b). Let  $C_{\Sigma}(\mathcal{A})$  be modular and  $a, x, y \in A$ ,  $a \neq x$ . Since  $\{x, y\} \subseteq C_{\mathcal{A}}(x) \vee C_{\mathcal{A}}(y)$ , we have

$$(*) \quad C_{\mathcal{A}}(x, y) \subseteq C_{\mathcal{A}}(x) \vee C_{\mathcal{A}}(y).$$

Suppose  $a \in C_{\mathcal{A}}(x, y)$ . Then  $a \in C_{\mathcal{A}}(x, y) \cap C_{\mathcal{A}}(a, x)$  and, by (\*), also

$$a \in (C_{\mathcal{A}}(x) \vee C_{\mathcal{A}}(y)) \cap C_{\mathcal{A}}(a, x).$$

Clearly  $C_{\mathcal{A}}(x) \subseteq C_{\mathcal{A}}(a, x)$  and, by modularity of  $C_{\Sigma}(\mathcal{A})$ , we conclude

$$a \in C_{\mathcal{A}}(x) \vee (C_{\mathcal{A}}(y) \cap C_{\mathcal{A}}(a, x)).$$

However,  $\mathcal{A}$  is  $\Sigma$ -separable, thus also  $a \in \{x\} \vee (\{y\} \cap C_{\mathcal{A}}(a, x))$ . Since  $a \neq x$ , this yields  $\{y\} \cap C_{\mathcal{A}}(a, x) \neq \emptyset$ , thus  $y \in C_{\mathcal{A}}(a, x)$  which proves (b).  $\square$

**Definition 2.** Let  $\varrho$  be a binary relation on  $A$ . We say that  $\varrho$  is *weakly transitive* if for each pairwise different elements  $a, b, c \in A$ ,  $\langle a, b \rangle \in \varrho$  and  $\langle b, c \rangle \in \varrho$  imply  $\langle c, a \rangle \notin \varrho$ .

**Corollary 1.** Let  $\mathcal{A} = (A, \{\varrho\})$  be an antisymmetrical binary relational system with one weakly transitive relation  $\varrho$  and  $C_\Sigma(\mathcal{A}) = \text{Conv } \mathcal{A}$ . The following conditions are equivalent:

- (1)  $\text{Conv } \alpha$  is modular;
- (2)  $\text{Conv } \mathcal{A}$  is distributive;
- (3) for any pairwise different elements  $a, b, c \in A$  we have  $\langle a, b \rangle \notin \varrho$  or  $\langle b, c \rangle \notin \varrho$ .

*Proof.* (3)  $\Rightarrow$  (2): If  $\mathcal{A}$  satisfies (3) then every subset of  $A$  is a convex subset, thus  $\text{Conv } \mathcal{A} = \text{Exp } A$ , i.e.  $\text{Conv } \mathcal{A}$  is distributive. (2)  $\Rightarrow$  (1) is trivial.

(1)  $\Rightarrow$  (3). Let  $\text{Conv } \mathcal{A}$  be modular and let  $a, x, y$  be pairwise different elements of  $A$ . Suppose  $x\varrho a$  and  $a\varrho y$ . Then  $a \in C_{\mathcal{A}}(x, y)$ ,  $a \notin x$  and  $y \notin C_{\mathcal{A}}(x, a)$  with respect to antisymmetry and weak transitivity of  $\varrho$ . Hence (b) of Theorem 2 is not valid. Moreover,  $\mathcal{A}$  is  $\Sigma$ -separable by Theorem 3 in [2] and, by Lemma 1,  $C_\Sigma(\mathcal{A})$  satisfies (i) of Theorem 1, thus we have a contradiction. Hence also (3) is satisfied.  $\square$

**Corollary 2.** Let  $\mathcal{S} = (S, \leq)$  be an ordered set and  $C_\Sigma(\mathcal{S}) = \text{Conv } \mathcal{S}$ . The following conditions are equivalent:

- (1)  $\text{Conv } \mathcal{S}$  is modular;
- (2)  $\text{Conv } \mathcal{S}$  is distributive;
- (3)  $\mathcal{S}$  does not contain a chain of length greater than two.

*Proof.* Clearly, any order is weakly transitive, and it is almost trivial to show that (3) of Corollary 1 is equivalent to (3) of Corollary 2 for  $\varrho = \leq$ .  $\square$

For any group  $\mathcal{G}$ , the lattice  $N(\mathcal{G})$  of all its normal subgroups is modular and it clearly satisfies (i) of Theorem 1 since  $\mathcal{G}_1 \vee \mathcal{G}_2 = \mathcal{G}_1 \cdot \mathcal{G}_2$  for each  $\mathcal{G}_1, \mathcal{G}_2 \in N(\mathcal{G})$ . However, it does not satisfy (ii) of Theorem 1: e.g. for the group  $(\mathbb{Z}, +)$  of all integers we have  $4 \in C_{\mathcal{G}}(2, 3) = \mathbb{Z}$ ,  $C_{\mathcal{G}}(4) \neq C_{\mathcal{G}}(2)$  but  $3 \notin C_{\mathcal{G}}(2, 4)$ . This motivates our effort to give another sufficient condition for modularity of  $C_\Sigma(\mathcal{A})$ . (Remark that a group  $\mathcal{G}$  is not  $\Sigma$ -separable with respect to  $C_\Sigma(\mathcal{G}) = N(\mathcal{G})$ .)

**Definition 3.** Let  $\mathcal{A} = (A, F, R)$  be an algebraic structure of type  $\tau$ . By a *binary formula* we mean any formula  $G(x_1, x_2, z, f)$  or  $G(x_1, x_2, z, \varrho)$  of the language  $L(\tau)$  provided  $f$  is a binary term operation of  $\mathcal{A}$  or  $\varrho$  is a binary relation of  $R$ .

**Theorem 3.** Let  $\mathcal{A} = (A, F, R)$  be an algebraic structure and let  $\Sigma$  contain a binary formula  $G(x_1, x_2, z, f)$  or  $G(x_1, x_2, z, \varrho)$  such that the following conditions are

satisfied:

(i) if  $X, Y \in C_\Sigma(\mathcal{A})$ ,  $\emptyset \neq X \neq Y \neq \emptyset$ , then  $a \in X \vee Y$  if and only if there exist  $b \in X, c \in Y$  such that  $G(b, c, a, f)$  or  $G(b, c, a, \varrho)$  is satisfied in  $\mathcal{A}$ ;

(ii) for each  $a, b, c \in A$ ,  $a \neq b$ , if the formula  $G(b, c, a, f)$  or  $G(b, c, a, \varrho)$  is satisfied in  $\mathcal{A}$  and  $C_{\mathcal{A}}(a) \neq C_{\mathcal{A}}(b)$  then  $c \in C_{\mathcal{A}}(a, b)$ . Then the lattice  $(C_\Sigma(\mathcal{A}), \subseteq)$  is modular.

*Proof.* Let  $X, Y, Z \in C_\Sigma(\mathcal{A})$  and  $X \subseteq Z$ . To check modularity of  $C_\Sigma(\mathcal{A})$  it is enough to consider  $\emptyset \neq X \neq Y \neq \emptyset$ . Suppose  $a \in (X \vee Y) \cap Z$ . By (i) there exist  $b \in X, c \in Y$  such that some binary formula  $G(b, c, a, f)$  or  $G(b, c, a, \varrho)$  is satisfied in  $\mathcal{A}$ . If  $C_\Sigma(a) = C_\Sigma(b)$  then

$$a \in C_{\mathcal{A}}(b) \subseteq X \subseteq X \vee (Y \cap Z).$$

If  $C_{\mathcal{A}}(a) \neq C_{\mathcal{A}}(b)$  then, by (ii),  $c \in C_{\mathcal{A}}(a, b)$ . However,  $a \in Z$  and  $b \in X \subseteq Z$ , thus also  $c \in C_{\mathcal{A}}(a, b) \subseteq Z$ . Hence, we conclude by (i)

$$a \in X \vee (Y \cap Z),$$

proving modularity of  $C_\Sigma(\mathcal{A})$ . □

*Example 3.* If  $\mathcal{G} = (A, \cdot, ^{-1}, e)$  is a group and  $C_\Sigma(\mathcal{G}) = N(\mathcal{G})$ , take a binary formula  $(x_1 \cdot x_2 = z)$ . Evidently, for  $\mathcal{A}_1, \mathcal{A}_2 \in N(\mathcal{G})$ ,  $a \in \mathcal{A}_1 \vee \mathcal{A}_2 = \mathcal{A}_1 \cdot \mathcal{A}_2$  if and only if there exist  $a_1 \in \mathcal{A}_1, a_2 \in \mathcal{A}_2$  with  $a = a_1 \cdot a_2$  and, if  $a = b \cdot c$  (i.e.  $G(b, c, a, \cdot)$  is satisfied in  $\mathcal{G}$ ) then  $c = b^{-1} \cdot a$ , thus  $c \in C_{\mathcal{A}}(a, b)$ . Hence, both (i), (ii) of Theorem 3 are satisfied.

*Example 4.* It is an easy exercise to verify that the quasiordered set of Example 1 also satisfies the assumptions of Theorem 3 for the binary formula  $G(x_1, x_2, z, \varrho) = (x_1 \varrho z \text{ and } z \varrho x_2)$ .

Now, we turn our attention to distributivity of  $C_\Sigma(\mathcal{A})$ .

**Theorem 4.** Let  $\mathcal{A} = (A, F, R)$  be an algebraic structure with the lattice  $C_\Sigma(\mathcal{A})$  of  $\Sigma$ -closed subsets. If there exists a binary term operation  $p(x, y)$  of  $\mathcal{A}$  such that

(i) for  $B, C \in C_\Sigma(\mathcal{A})$  we have  $a \in B \vee C$  if and only if  $a = p(b, c)$  for some  $b \in B, c \in C$ ;

(ii) if  $D \in C_\Sigma(\mathcal{A})$  and  $p(b, c) \in D$  for some  $b, c \in A$ , then  $b, c \in D$ , then the lattice  $(C_\Sigma(\mathcal{A}), \subseteq)$  is distributive.

*Proof.* Suppose  $B, C, D \in C_\Sigma(\mathcal{A})$  and  $a \in D \cap (B \vee C)$ . Then  $a \in D$  and, by (i), there exist  $b \in B, c \in C$  with  $a = p(b, c)$ . Hence also  $p(b, c) \in D$  and, by (ii), we have  $b \in D, c \in D$ . Thus  $b \in D \cap B, c \in D \cap C$  and by (i) again, we conclude  $a = p(b, c) \in (D \cap B) \vee (D \cap C)$ . □

**Example 5.** If  $\mathcal{L}$  is a distributive lattice and  $C_\Sigma(\mathcal{L}) = \text{Id } \mathcal{L}$ , we can put  $p(x, y) = x \vee y$ . It is well-known that for  $I_1, I_2 \in \text{Id } \mathcal{L}$ ,  $y \in I_1 \vee I_2$  if and only if  $y = i_1 \vee i_2$  for some  $i_1 \in I_1, i_2 \in I_2$ . Moreover, if  $J \in \text{Id } \mathcal{L}$  and  $j_1 \vee j_2 \in J$  for  $j_1, j_2 \in \mathcal{L}$  then  $j_1 \leq j_1 \vee j_2, j_2 \leq j_1 \vee j_2$  imply also  $j_1, j_2 \in J$ .

Thus both assumptions of Theorem 4 are satisfied.

Now, let  $\mathcal{A} = (A, F, R)$  be an algebraic structure and let  $B \in C_\Sigma(\mathcal{A})$  for some given set  $\Sigma$  of open formulas. If there exists an element  $b \in A$  such that  $B = C_{\mathcal{A}}(b)$ , we say that  $b$  is a *generator* of  $B$ .

In the remaining part of the paper, denote by  $Z$  the set of all integers and suppose  $F \neq \emptyset$  for any algebraic structure  $\mathcal{A} = (A, F, R)$  under consideration.

**Definition 4.** An algebraic structure  $\mathcal{A} = (A, F, R)$  is called  $\Sigma$ -cyclic if there exist an element  $d \in A$ , a subset  $K \subseteq Z$  and binary integral operations  $\varphi, \psi: K \times K \rightarrow K$  and unary terms  $w_k(x)$  for  $k \in K$  of  $\mathcal{A}$  such that

- (a) for each  $B \in C_\Sigma(\mathcal{A})$  there exists  $k \in K$  such that  $w_k(d)$  is a generator of  $B$ ;
- (b) if  $w_m(d)$  or  $w_n(d)$  are generators of  $B$  or  $D$ , respectively, for  $B, D \in C_\Sigma(\mathcal{A})$ , then  $w_{\varphi(m,n)}(d)$  or  $w_{\psi(m,n)}(d)$  are generators of  $B \vee D$  or  $B \cap D$ , respectively;
- (c)  $\psi(k, \varphi(m, n)) = \varphi(\psi(k, m), \psi(k, n))$  for every  $k, m, n \in K$ .

The terms  $w_k(x)$  are called *characteristic terms* of  $C_\Sigma(\mathcal{A})$ .

**Theorem 5.** If  $\mathcal{A} = (A, F, R)$  is a  $\Sigma$ -cyclic algebraic structure then the lattice  $(C_\Sigma(\mathcal{A}), \subseteq)$  is distributive.

**Proof.** Let  $\mathcal{A}$  be a  $\Sigma$ -cyclic algebraic structure and let  $w_k(x)$  be its characteristic terms for  $k \in K \subseteq Z$ . Suppose that  $\varphi$  and  $\psi$  satisfy (b) and (c) of Definition 4. Let  $B, C, D \in C_\Sigma(\mathcal{A})$ . Suppose that  $d \in A$  and  $w_m(d)$  or  $w_n(d)$  or  $w_k(d)$  are generators of  $B$  or  $C$  or  $D$ , respectively. By (a), (b), (c) of Definition 4, we can easily derive

$$\begin{aligned} D \cap (B \vee C) &= C_{\mathcal{A}}(w_k(d)) \cap (C_{\mathcal{A}}(w_n(d)) \vee C_{\mathcal{A}}(w_m(d))) \\ &= C_{\mathcal{A}}(w_{\psi(k, \varphi(m, n))}(d)) = C_{\mathcal{A}}(w_{\varphi(\psi(k, m), \psi(k, n))}(d)) \\ &= (C_{\mathcal{A}}(w_k(d)) \cap C_{\mathcal{A}}(w_m(d))) \vee (C_{\mathcal{A}}(w_k(d)) \cap C_{\mathcal{A}}(w_n(d))) \\ &= (D \cap B) \vee (D \cap C), \end{aligned}$$

i.e. the lattice  $(C_\Sigma(\mathcal{A}), \subseteq)$  is distributive.  $\square$

**Example 6.** If  $\mathcal{G} = (G, \cdot)$  is a cyclic group and  $C_\Sigma(\mathcal{G}) = \text{Sub } \mathcal{G}$ , put  $K = Z$ ,  $w_k = x^k$  and  $\varphi(m, n) = \text{GCD}(m, n)$ ,  $\psi(m, n) = \text{LCM}(m, n)$ . As an element  $d \in G$  we pick up the generator of  $\mathcal{G}$ . Evidently,  $\mathcal{G}$  is  $\Sigma$ -cyclic.



Example 7. If  $\mathcal{A} = (A, f)$  is a monounary algebra and  $C_\Sigma(\mathcal{A}) = \text{Sub } \mathcal{A}$ , we can put  $K = \mathbb{N} \cup \{0\}$  (non-negative integers),  $w_k(x) = f^k(x)$  where  $f^0(x) = x$  and  $f^{k+1}(x) = f(f^k(x))$  for each  $k \in K$ . Moreover, put  $\varphi(m, n) = \min(m, n)$ ,  $\psi(m, n) = \max(m, n)$ . If  $\mathcal{A}$  has a unique generator  $d$  then  $\mathcal{A}$  is  $\Sigma$ -cyclic.

Example 8. Suppose  $\mathcal{A} = (A, F, R)$  is an algebraic structure with at least two elements such that  $F$  contains a nullary operation  $c$  and  $f(c, \dots, c) = c$  for each  $f \in F$ . Further, suppose  $C_\Sigma(\mathcal{A}) = \{\{c\}, A\}$  (trivially,  $C_\Sigma(\mathcal{A})$  is distributive). Put  $K = \{0, 1\}$ . If  $A \neq \{c\}$ , choose  $d \neq c$ ,  $d \in A$  and put  $w_0(x) = c$ ,  $w_1(x) = d$ . Further, let  $\varphi$  and  $\psi$  be defined in the same manner as in the foregoing Example 7. Evidently,  $\mathcal{A}$  is  $\Sigma$ -cyclic.

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