

Descriptive complexity of Banach spaces

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Why?

- let P be a property of sBSs invariant under isomorphism / isometry
- then we may talk about the **set** of codes of those sBSs having property P
- we may investigate the **complexity** of this set (Borel, analytic, ... / open, G_δ , ...)

Theorem (Szlenk, 1968; Bossard, 1993)

There is no reflexive sBS containing an isomorphic copy of every reflexive sBS.

Proof: In a certain coding of sBSs it holds:

- the set of codes of all subspaces of a fixed sBS is analytic
- the set of codes of all reflexive sBSs is not analytic (Bossard, 1993)



The 'classical' coding of sBSs by a standard Borel space

$C([0, 1])$ contains an isometric copy of every sBS

$\mathcal{F}(C([0, 1])) \equiv$ the set of all closed subsets of $C([0, 1])$

$\mathcal{F}(C([0, 1]))$ is equipped with the **Effros-Borel structure**, that is, with the σ -algebra generated by the sets

$$\{F \in \mathcal{F}(C([0, 1])) : F \cap U \neq \emptyset\}, \quad U \subseteq C([0, 1]) \text{ open}$$

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Fact

$\mathcal{F}(C([0, 1]))$ is a standard Borel space.

Fact

$SB := \{F \in \mathcal{F}(C([0, 1])) : F \text{ is linear}\}$ is a Borel subset of $\mathcal{F}(C([0, 1]))$. In particular, it is a standard Borel space.

Note that there is no canonical Polish topology on SB compatible with the Effros-Borel structure.

Coding of sBSs by a Polish space (Godefroy, Saint-Raymond, 2018)

A Polish topology on $\mathcal{F}(C([0, 1]))$ is called **admissible** if it satisfies certain natural axioms.

Properties of admissible topologies:

- The set $SB := \{F \in \mathcal{F}(C([0, 1])) : F \text{ is linear}\}$ is a G_δ -set. In particular, it is a Polish space.
- The Borel σ -algebra generated by an admissible topology is the Effros-Borel structure.
- For two admissible topologies, the identity function is of Baire class 1.

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Examples of admissible topologies on $\mathcal{F}(C([0, 1]))$

- the Vietoris topology 'inherited' from $F(\hat{C}([0, 1]))$, where a compatible totally bounded metric on $C([0, 1])$ is fixed and $\hat{C}([0, 1])$ is the corresponding metric completion
- the Wijsman topology generated by the maps $F \in \mathcal{F}(C([0, 1])) \mapsto d(F, f)$, $f \in C([0, 1])$, where d is a compatible metric on $C([0, 1])$

Coding of sBSs by the Polish space of (pseudo)norms

Let V be the vector space over \mathbb{Q} of all finitely supported sequences of rational numbers.

Definition

Let $\mathcal{P} \subseteq \mathbb{R}^V$ be the set of all pseudonorms on V .

Then $\mu \in \mathcal{P}$ is a code for the sBS X_μ obtained as follows:

1. extend μ to a pseudonorm on c_{00} (over \mathbb{R})
2. take the quotient $(c_{00}, \mu) / \{x \in c_{00} : \mu(x) = 0\}$
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Fact

- $\mathcal{P} \subseteq \mathbb{R}^V$ is a closed set
- $\mathcal{P}_\infty := \{\mu \in \mathcal{P} : X_\mu \text{ is infinite-dimensional}\}$ is a G_δ -set
- $\mathcal{B} := \{\mu \in \mathcal{P} : \mu \text{ is a norm on } c_{00}\}$ is a G_δ -set

Fact

For every infinite-dimensional sBS X there is $\mu \in \mathcal{B}$ such that $X \equiv X_\mu$.

Proof: Let f_1, f_2, \dots be linearly independent vectors in X such that $\overline{\text{span}}(f_1, f_2, \dots) = X$.

Define $\mu: V \rightarrow \mathbb{R}$ by

$$\mu \left(\sum_{i \in F} \alpha_i e_i \right) := \left\| \sum_{i \in F} \alpha_i f_i \right\|_X.$$



Three ways of coding infinite-dimensional sBSs by a Polish space:
(a) SB_∞ with an admissible topology, (b) \mathcal{P}_∞ , (c) \mathcal{B} .

Theorem (Informal statement)

For a given class of Banach spaces, the Borel complexity of the set of the corresponding codes depends only a little on the choice of the coding.

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Theorem (Formal statement)

- *There is a continuous map $\varphi: SB_\infty \rightarrow \mathcal{P}_\infty$ such that $Y \equiv X_{\varphi(Y)}$, $Y \in SB_\infty$.*
- *There is a Baire class 1 map $\psi: \mathcal{P}_\infty \rightarrow \mathcal{B}$ such that $X_\mu \equiv X_{\psi(\mu)}$, $\mu \in \mathcal{P}_\infty$.*
- *There is a Baire class 1 map $\chi: \mathcal{B} \rightarrow SB_\infty$ such that $X_\mu \equiv \chi(\mu)$, $\mu \in \mathcal{B}$.*

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From now on we use the coding \mathcal{B} only.

Definition

Let X, Y be Banach spaces. We say that X is **finitely representable** in Y if for every finite-dimensional subspace $E \subseteq X$ and every $\varepsilon > 0$ there exists an isomorphic embedding $T : E \rightarrow Y$ such that $\|T\| \|T^{-1}\| \leq 1 + \varepsilon$.

For a sBS X we denote by $\langle X \rangle_{\equiv}^{\mathcal{B}}$ the set $\{\mu \in \mathcal{B} : X_{\mu} \equiv X\}$.

Fact

Let X be a sBS. Then

$$\overline{\langle X \rangle_{\equiv}^{\mathcal{B}}} = \{\mu \in \mathcal{B} : X_{\mu} \text{ is finitely representable in } X\}.$$

Complexity of isometry classes and isomorphism classes

For a sBS X we denote

- $\langle X \rangle_{\equiv}^{\mathcal{B}} = \{\mu \in \mathcal{B} : X_{\mu} \equiv X\}$
- $\langle X \rangle_{\simeq}^{\mathcal{B}} = \{\mu \in \mathcal{B} : X_{\mu} \simeq X\}$

Fact

Let X be a sBS. Then $\langle X \rangle_{\equiv}^{\mathcal{B}}$ is a Borel set.

Fact

Let X be a sBS. Then $\langle X \rangle_{\simeq}^{\mathcal{B}}$ is an analytic set but not necessarily Borel.

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Fact

Let X be a sBS. Then $\langle X \rangle_{\equiv}^{\mathcal{B}}$ is a Borel set.

Proof:

- the relation of linear isometry is Borel bireducible with an orbit equivalence relation (Melleray, 2007)
- orbit equivalence relations have Borel equivalence classes □

Fact

Let X be a sBS. Then $\langle X \rangle_{\simeq}^{\mathcal{B}}$ is an analytic set but not necessarily Borel.

Proof: These spaces do not have a Borel isomorphism class:

$C(2^{\omega})$, $L_p([0, 1])$ ($1 < p < \infty, p \neq 2$), c_0 , ...



Theorem

ℓ_2 is the unique, up to isometry, infinite-dimensional sBS X such that $\langle X \rangle_{\equiv}^{\mathcal{B}}$ is a closed set.

Note that $\langle X \rangle_{\equiv}^{\mathcal{B}}$ is never an open set.

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Proof: ℓ_2 is characterized by the parallelogram law

$$2\|x\|^2 + 2\|y\|^2 = \|x + y\|^2 + \|x - y\|^2, \quad x, y \in \ell_2.$$

Thus $\langle \ell_2 \rangle_{\equiv}^{\mathcal{B}}$ is a closed set.

Now suppose that $\langle X \rangle_{\equiv}^{\mathcal{B}}$ is a closed set.

By Dvoretzky's theorem, ℓ_2 is finitely representable in X .

So $\overline{\langle X \rangle_{\equiv}^{\mathcal{B}}}$ contains all $\mu \in \mathcal{B}$ for which $X_{\mu} \equiv \ell_2$.

On the other hand, all elements of $\langle X \rangle_{\equiv}^{\mathcal{B}} = \overline{\langle X \rangle_{\equiv}^{\mathcal{B}}}$ are codes for spaces isometric to X . □

Theorem

ℓ_2 is the unique, up to isomorphism, infinite-dimensional sBS X such that $\langle X \rangle_{\cong}^{\mathcal{B}}$ is an F_σ set.

Note that $\langle X \rangle_{\cong}^{\mathcal{B}}$ is never an open/closed set.

We do not know whether $\langle X \rangle_{\cong}^{\mathcal{B}}$ can be a G_δ set but the only candidate is the Gurariĭ space.

Theorem

ℓ_2 is the unique, up to isomorphism, infinite-dimensional sBS X such that $\langle X \rangle_{\simeq}^{\mathcal{B}}$ is an F_σ set.

Note that $\langle X \rangle_{\simeq}^{\mathcal{B}}$ is never an open/closed set.

We do not know whether $\langle X \rangle_{\simeq}^{\mathcal{B}}$ can be a G_δ set but the only candidate is the Gurariĭ space.

Proof: By Kwapien's theorem, ℓ_2 is the unique (up to isomorphism) sBS that has type 2

$(\exists c > 0 \forall x_1, \dots, x_n \in X : (\mathbb{E} \|\sum_{i=1}^n \pm x_i\|^2)^{1/2} \leq c (\sum_{i=1}^n \|x_i\|^2)^{1/2})$ and cotype 2

$(\exists c > 0 \forall x_1, \dots, x_n \in X : (\sum_{i=1}^n \|x_i\|^2)^{1/2} \leq c (\mathbb{E} \|\sum_{i=1}^n \pm x_i\|^2)^{1/2})$.

'Having type/cotype 2' are F_σ conditions.

So $\langle \ell_2 \rangle_{\simeq}^{\mathcal{B}}$ is an F_σ set.

The other implication is based on the solution to the homogeneous subspace problem (Komorowski, Tomczak-Jaegermann, 1995; Gowers, 2002).

Definition

The Gurariĭ space is the unique, up to isometry, sBS \mathbb{G} such that for every $\varepsilon > 0$, every finite-dimensional Banach spaces $A \subseteq B$ and every isometric embedding $e: A \rightarrow \mathbb{G}$ there is an extension $f: B \rightarrow \mathbb{G}$ of e such that f is an ε -isometric embedding.

Theorem

The isometry class $\langle \mathbb{G} \rangle_{\cong}^{\mathcal{B}}$ is a dense G_δ set.

We do not know the descriptive complexity of the isomorphism class $\langle \mathbb{G} \rangle_{\simeq}^{\mathcal{B}}$.

The isometry class of L_p

Theorem

For every $1 \leq p < \infty$, $p \neq 2$, the isometry class $\langle L_p([0, 1]) \rangle_{\cong}^{\mathcal{B}}$ is a G_δ -complete set.

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Proof:

- The class of $\mathcal{L}_{p,1+}$ sBSs is a G_δ set (A sBS X is called an $\mathcal{L}_{p,1+}$ -space if for every finite-dimensional $E \subseteq X$ and $\varepsilon > 0$ there is a finite-dimensional $E \subseteq F \subseteq X$ and a linear isomorphism $T : F \rightarrow \ell_p^n$ with $\|T\| \|T^{-1}\| \leq 1 + \varepsilon$.)
- An $\mathcal{L}_{p,1+}$ sBS X is isometric to $L_p([0, 1])$ if and only if

$$\forall x \in S_X \quad \forall \varepsilon > 0 \quad \forall \delta > 0 \quad \exists x_1, x_2 \in X :$$

$$(x_1, x_2) \stackrel{1+\varepsilon}{\sim} \ell_p^2 \quad \text{and} \quad \|2^{1/p}x - x_1 - x_2\| < \delta.$$



The isometry class of ℓ_p

Theorem

For every $1 \leq p < \infty$, $p \neq 2$, the isometry class $\langle \ell_p \rangle_{\equiv}^{\mathcal{B}}$ is an $F_{\sigma\delta}$ -complete set.

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Proof:

- The class of $\mathcal{L}_{p,1+}$ sBSs is a G_δ set.
- An $\mathcal{L}_{p,1+}$ sBS X is isometric to ℓ_p if and only if

$\forall x \in S_X \forall \delta \in (0, 1) \exists \varepsilon > 0 \exists N \in \mathbb{N} \forall x_1, \dots, x_N \in X :$

$$(N^{1/p} x_i)_{i=1}^N \stackrel{1+\varepsilon}{\sim} \ell_p^N \Rightarrow \|x - \sum_{i=1}^N x_i\| > \delta.$$



The isometry class of c_0

Theorem

The isometry class $\langle c_0 \rangle_{\equiv}^{\mathcal{B}}$ is an $F_{\sigma\delta}$ -complete set.

Theorem

The isometry class $\langle c_0 \rangle_{\cong}^{\mathcal{B}}$ is an $F_{\sigma\delta}$ -complete set.

Proof:

- The class of $\mathcal{L}_{\infty,1+}$ sBSs is a G_{δ} set (A sBS X is called an $\mathcal{L}_{\infty,1+}$ -space if for every finite-dimensional $E \subseteq X$ and $\varepsilon > 0$ there is a finite-dimensional $E \subseteq F \subseteq X$ and a linear isomorphism $T : F \rightarrow \ell_{\infty}^n$ with $\|T\| \|T^{-1}\| \leq 1 + \varepsilon$.)
- Let $0 < \varepsilon < 1$. An $\mathcal{L}_{\infty,1+}$ sBS X is isometric to c_0 if and only if

$$(B_{X^*})'_{2\varepsilon} = (1 - \varepsilon)B_{X^*},$$

where $F'_{\varepsilon} = \{x^* \in F : U \ni x^* \text{ is } w^*\text{-open} \Rightarrow \text{diam}(U \cap F) \geq \varepsilon\}$ is the Szlenk derivative.

- Let $\varepsilon > 0$. The mapping

$$F \mapsto F'_{\varepsilon}, \quad F \subset X^* \text{ is } w^*\text{-compact},$$

is of Baire class 2.



M. Cúth, M. Doležal, M. Doucha, O. Kurka.
Polish spaces of Banach spaces.
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